

Chapter 5

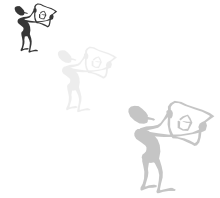
Computer Organization

Process Control

Outline

➔ ● **General Architecture of Computer**

- CPU - MEM - I/O
- Peripheral

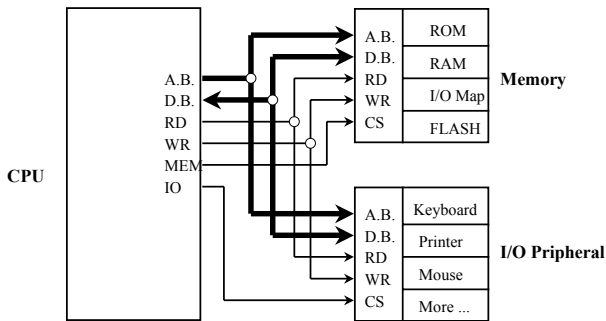


● **Memory Hierarchy**

- Main Memory
- Auxiliary Memory

● **Input-Output Interface**

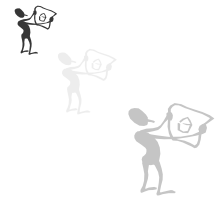
CPU - Memory - I/O



Outline

● **General Architecture of Computer**

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- Peripheral

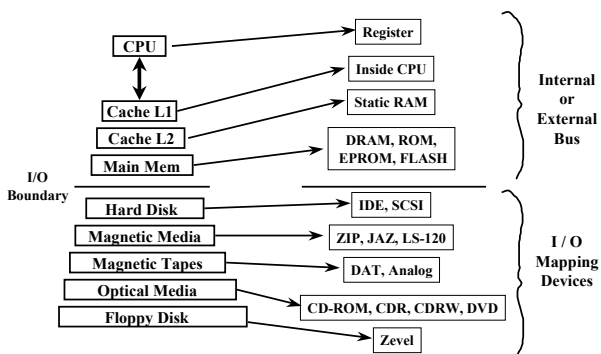


➔ ● **Memory Hierarchy**

- Main Memory
- Auxiliary Memory

● **Input-Output Interface**

Memory Hierarchy



Main Memory

● **RAM (Random Access Memory):**

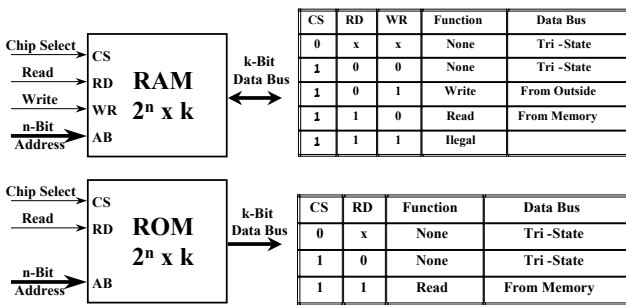
- *Static* => Flip Flop, Fast, Expensive, No Controller.
- *Dynamic* => Mos Capacitor, Slow, Cheap, Refreshing controller.

● **ROM (Read Only Memory):**

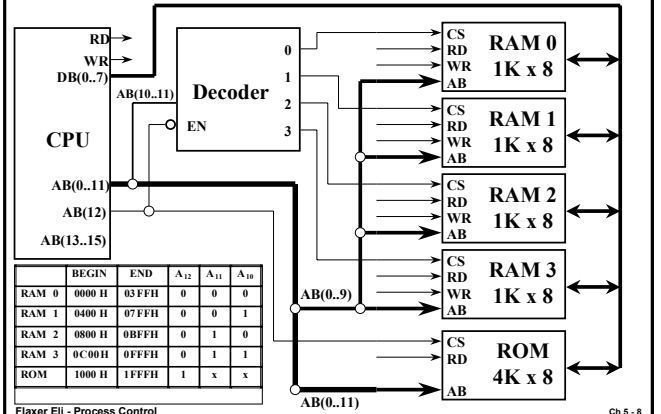
- Type: ROM, PROM, EPROM, E²PROM, Flash.
- Used: BIOS, Table, etc.

- Not all the memory address space must be occupied, parts of them can be empty (without memory or other devices).

RAM & ROM Chips



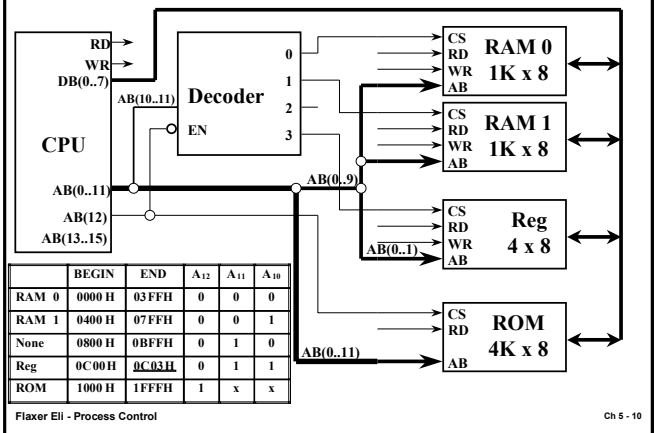
Memory Address Map



Address Map Table

	BEGIN	END	A ₁₂	A ₁₁	A ₁₀
RAM 0	0000 H	03 FFH	0	0	0
RAM 1	0400 H	07 FFH	0	0	1
RAM 2	0800 H	0B FFH	0	1	0
RAM 3	0C00 H	0F FFH	0	1	1
ROM	1000 H	1F FFH	1	x	x

Unoccupied Memory Address



Address Map Table

	BEGIN	END	A ₁₂	A ₁₁	A ₁₀
RAM 0	0000 H	03 FFH	0	0	0
RAM 1	0400 H	07 FFH	0	0	1
None	0800 H	0B FFH	0	1	0
Reg	0C00 H	0C03 H	0	1	1
ROM	1000 H	1F FFH	1	x	x

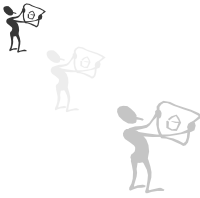
Auxiliary Memory

- **Hard Disk:**
 - IDE => 30G, Fast, Internal.
 - Scsi => 100G, Ultra Fast, Internal / External.
- **TAPE :**
 - DAT: Slow, 100G per cassette.
 - Travan: Very Slow, 1G per cassette.
- **Optics:**
 - CD, CDR, CDRW => 640M.
 - DVD, DVD-RAM => 5G.
- **Removable:**
 - Zip => 100M, 250M.
 - Jaz => 1G, 2G.
 - Ls-120 => 120M.
 - More => Floppy, Flash, ...

Outline

- **General Architecture of Computer**

- CPU - MEM - I/O
- Peripheral



- **Memory Hierarchy**

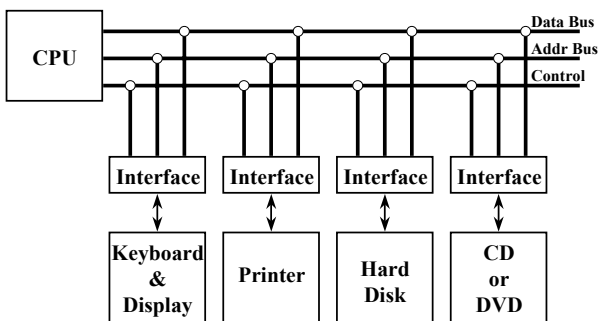
- Main Memory
- Auxiliary Memory

- ➔ ● **Input-Output Interface**

Input-Output Interface

- I/O interface provide a method for transferring information between internal storage and external I/O devices.
- To resolve the differences between CPU and peripherals (Mechanism, Transfer Rate, Format, etc) the system include interface.
- In addition, each device may have its own controller that supervise the operation of the particular mechanism in the peripheral.

I/O Bus and Interface



Input-Output Interface

- To communicate with particular device, the processor place a device address on the address bus (**Port Addr**).
- When the interface detects its own address, it activates the device that it controls. All others peripherals are disable in Tri-State.
- At the same time the processor provides a function code in the control lines.
- The data send / received in Bi-directional data bus.
- Each CPU has special opcode for I/O operation.

Physical Memory Accessing

Assembler	READ	WRITE
	Mov DI, AbsAdr	Mov DI, AbsAdr
	Mov AI, [DI]	Mov [DI], AI
C & C++	READ	WRITE
	char *AbsAdrPtr = AbsAdr;	
	MyData = (*AbsAdrPtr);	(*AbsAdrPtr) = MyData;

Physical Memory Accessing

- If the operation system not allow direct accessing to the hardware, we must use tool kit or lib to map the physical memory to the pointer.
- Each compiler has its own library

```
static unsigned char *kbf1;
int pt1;
int main ()
{
    unsigned char temp;
    MapPhysicalMemory (0x417, 1, &kbf1, &pt1);
    temp = *kbf1;
    printf("%x", temp);
    UnMapPhysicalMemory (pt1);
}
```

Physical Memory Accessing

- ReadFromPhysicalMemory (AbsAdr, Buffer, Number_Of_Byte);
- WriteToPhysicalMemory (AbsAdr, Buffer, Number_Of_Byte);

```
int main ()
{
    unsigned char temp;
    ReadFromPhysicalMemory (0x417, temp, 1);
    WriteToPhysicalMemory (0x469, 1000, 4);
    printf("%x", temp);
}
```

Input-Output Instruction

Assembler	IN	OUT
BYTE	In AL, Op1	Out Op1, AL
WORD	In AX, Op1	Out Op1, AX
DWORD	In EAX, Op1	Out Op1, EAX

Op1 = Immd 8 bit Addr or 16 bit address in DX.

C & C++	IN	OUT
BYTE char	inp (int Address)	outp (int Address, char Data)
WORD short	inpw (int Address)	outpw (int Address, short Data)

Input-Output Instruction

PASCAL & DELPHI 1.0

BYTE Data := Port[Address]	Port[Address] := Data
WORD Data := Portw[Address]	Portw[Address] := Data

DELPHI 2.0+ (32 bit)

Not Support	Not Support
Use inline assembler	Use inline assembler

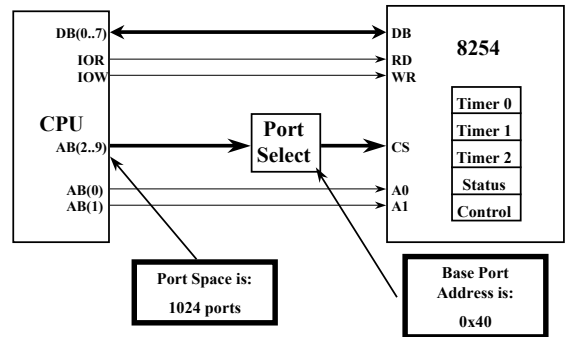
BASIC

Data% = INP(Address%)	OUT(Address%, Data%)
-----------------------	----------------------

VISUAL BASIC

Not Support	Not Support
Use external DLL	Use external DLL

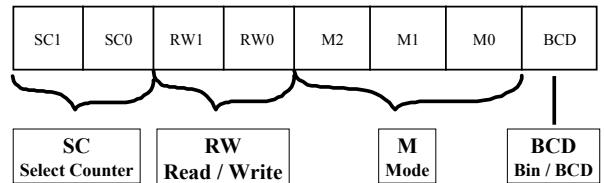
Example: 8254



8254 Addresses

Select	CS	A ₁	A ₀	RD	WR
Tri -State	0	x	x	x	x
Timer 0	1	0	0	x	x
Timer 1	1	0	1	x	x
Timer 2	1	1	0	x	x
Control	1	1	1	0	1
Status	1	1	1	1	0

8254 Control



8254 Control Field

SC - Select Counter

SC1	SC0	
0	0	Select Counter 0
0	1	Select Counter 1
1	0	Select Counter 2
1	1	Read - Back Command

RW - Read Write

RW1	RW0	
0	0	Counter Latch Command
0	1	Read / Write LSB
1	0	Read / Write MSB
1	1	Read / Write LSB first & MSB after

8254 Control Field

M - Mode

M2	M1	M0	
0	0	0	Mode 0
0	0	1	Mode 1
0	1	0	Mode 2
0	1	1	Mode 3
1	0	0	Mode 4
1	0	1	Mode 5

BCD - Binary Code Decimal

BCD	
0	Binary Counter 16 Bit
1	BCD Counter 4 Digit

I/O Programming Example (C)

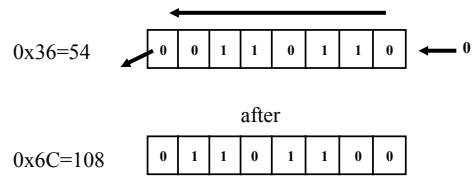
- Set the Timer2 of 8254 to mode 3, Binary counting, and 16 bit interfacing. Load the counter with 0x1234.

1 0 1 1 0 1 1 0 = 0xB6

```
int main ()
{
    unsigned char temp;
    temp = 0xB6;
    outp(0x43, temp) // write control
    outp(0x42, 0x34) // write least
    outp(0x42, 0x12) // write most
}
```

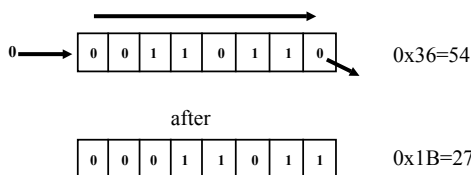
Shift Operations

- C language has two shift operations for variable and constant:
 - >> (shift right) <<< (shift left).
- Shift left operation always insert '0' to the LSB



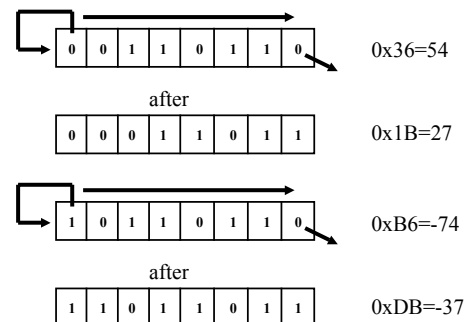
Shift Operations

- Shift right operation insert, to the MSB, value that depended in the type of the variable or the constant.
- If the type is unsigned, '0' is insert to the MSB



Shift Operations

- If the type is sign, the MSB is duplicated (sign extended).



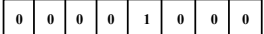
Shift Operations Example (C)

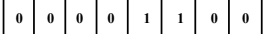
```
int main ()
{
    unsigned char x, y;
    char a, b;
    x = 1;           // x = 00000001
    y = x << 3;      // y = 8
    x = 255;        // x = 11111111
    y = x >> 3;      // y = 31

    a = 1;          // a = 00000001
    b = a << 3;      // b = 8
    a = -1;         // a = 11111111
    b = a >> 3;      // b = -1
}
```

Constant Shift Operations

- The shift operations work with a constant too.
- For example:

— $1 \ll 3 = 8$ 

— $48 \gg 2 = 12$ 

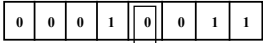
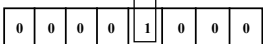
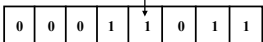
- Constant integer, is by default 32 bit signed.
- Unsigned constant is defined by U.

Bit Manipulations

- How to set especial bit of the variable ?
- Set the k bit:

— $x |= (1 \ll k)$

— For example set bit 3 $X |= (1 \ll 3)$:

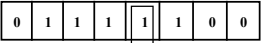
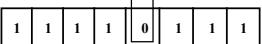
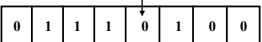
	X = 0x13
	1 << 3
	X = 0x1B

Bit Manipulations

- How to clear especial bit of the variable ?
- Clear the k bit:

— $x \&= \sim(1 \ll k)$

— For example clear bit 3 $X \&= \sim(1 \ll 3)$:

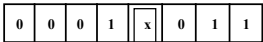
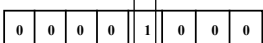
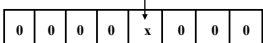
	X = 0x7C
	$\sim(1 \ll 3)$
	X = 0x74

Bit Sensing

- How to check especial bit of the variable ?
- Check the k bit of variable Y:

— $Y \& (1 \ll k)$

— For example check bit 3 $Y \& (1 \ll 3)$:

	Y
	1 << 3
	

If (x == 0) the all expression is FALSE
If (x == 1) the all expression is TRUE

Byte to Bit Conversion Example (C)

```
Void Byte2Bit (char b, char y[])
{
    int k;
    for (k=0; k<8; k++)
        y[k] = b & (1<<k);
}
```

```
Void Byte2Bit (char b, char y[])
{
    int k;
    for (k=0; k<8; k++, b >>= 1)
        y[k] = b & 1;
}
```

Bit to Byte Conversion Example (C)

```
char Bit2Byte (char y[])
{
    int k, b;
    for (k=0, b=0; k<8; k++)
        b |= (y[k] << k);
    return(b);
}
```

```
char Bit2Byte (char y[])
{
    int k, b;
    for (k=7, b=0; k>=0; b |= y[k--])
        b <<= 1;
    return(b);
}
```